Quantum graphs and their infinite path spaces

Lara Ismert

joint w/ M. Brannan, M. Hamidi, B. Nelson, M. Wasilewski

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A story about classical graphs

Let G = (V, A) be a finite simple graph.

▶ $\{0,1\}$ -matrix $A \rightsquigarrow$ Cuntz–Krieger algebra \mathcal{O}_A

- edge correspondence C(E) (over C(V)) $\rightsquigarrow \mathcal{O}_{C(E)}$
- ► Exel system $(C(E^{\infty}), \alpha, \mathcal{L}) \rightsquigarrow C(E^{\infty}) \rtimes_{\alpha, \mathcal{L}} \mathbb{N}$

"Fact":
$$\mathcal{O}_A \cong \mathcal{O}_{\mathcal{C}(E)} \cong \mathcal{C}(X_A) \rtimes_{\alpha, \mathcal{L}} \mathbb{N}.$$

A classical graph's Exel crossed product

Let
$$G = (V, A)$$
 be a finite simple graph. Define

$$E^{\infty} := \{ (v_{k_1}, v_{k_2}, ...) : v_{k_i} \in V, A_{k_i k_{i+1}} = 1 \ \forall i \in \mathbb{N} \}.$$

Consider the left shift $\sigma: E^\infty \to E^\infty$ given by

$$\sigma(v_{k_1}, v_{k_2}, ...) := (v_{k_2}, v_{k_3}, ...).$$

Definition

If G has no sinks, define $\alpha, \mathcal{L}: C(E^{\infty}) \rightarrow C(E^{\infty})$ by:

$$\begin{array}{l} \bullet \ \alpha(f) := f \circ \sigma \\ \bullet \ [\mathcal{L}(f)](\theta) := \frac{1}{|\sigma^{-1}(\{\theta\})|} \sum_{\gamma \in \sigma^{-1}(\{\theta\})} f(\gamma) \quad \forall \theta \in E^{\infty} \end{array}$$

 $(C(E^{\infty}), \alpha, \mathcal{L}) \quad \rightsquigarrow \quad \text{Exel crossed product } C(E^{\infty}) \rtimes_{\alpha, \mathcal{L}} \mathbb{N}.$

In a quantum graph...

Definition

a finite vertex set is replaced by a quantum set consisting of

- ▶ a (finite-dimensional) C*-algebra B
- a state ψ that is a δ -form

Definition

The adjacency matrix is replaced by a *quantum adjacency matrix*, a linear map $A: B \rightarrow B$ which is **quantum idempotent**:

$$m(A\otimes A)m^* = \delta^2 A$$

Example

Let (V, A) be a finite graph. Then B := C(V) is a f.d. C^* -algebra. Let $\{p_v : v \in V\}$ be the basis for C(V), and define

$$\psi(p_{m{v}}):=rac{1}{|m{V}|} \ \ orall m{v}\inm{V}.$$

 $(C(V), \psi)$ is a quantum set and $(C(V), \psi, A)$ is a quantum graph.

Remark: All finite simple graphs can be viewed as quantum graphs. Note: $C(E^2) \cong C(E) \otimes_{C(V)} C(E)$, and $C(E^{\infty}) \cong C(E)^{\otimes_{C(V)} \mathbb{N}}$

Let (B, ψ) be a finite quantum set.

Example (Complete quantum graph) Define $A : B \rightarrow B$ on $x \in B$ by

$$A(x) = \delta^2 \psi(x) \mathbf{1}_B.$$

Notation: $K(B, \psi)$

Edges and Paths

- $E \cong B \otimes_{\psi} B$ ("all possible pairs of vertices")
- $\blacktriangleright \ E^2 \cong E \otimes_B E \cong (B \otimes_{\psi} B) \otimes_B (B \otimes_{\psi} B) \cong B \otimes_{\psi} B \otimes_{\psi} B$

Let (B, ψ) be a finite quantum set.

Example (Trivial quantum graph) Define $A : B \rightarrow B$ on $x \in B$ by

$$A(x) = x.$$

Notation: $T(B, \psi)$

Edges and Paths

- $E \cong B$ ("just a loop at each vertix")
- $\blacktriangleright E^2 \cong E \otimes_B E \cong B \otimes_B B \cong B$

Let \mathcal{G} be a quantum graph.

- BEVW '21 and BHINW '22 defined a quantum Cuntz-Krieger algebra O(G)
- BHINW '22 defined a quantum analogue of C(E), from which we can build a Cuntz-Pimsner algebra.
- We have examples of a reasonable is the infinite path space for a quantum graph?
- Is there a natural Exel system whose crossed product recovers the quantum graph's Cuntz-Krieger algebra?

Exel systems for some quantum graphs

Example (Complete quantum graph)

The quantum edge correspondence is $E \cong B \otimes_{\psi} B$

 \rightsquigarrow quantum infinite path space should contain "all possible paths."

Consider $E^{\otimes_B \mathbb{N}} = (B \otimes_{\psi} B) \otimes_B (B \otimes_{\psi} B) \dots \cong B^{\otimes \mathbb{N}}$ with $\alpha, \mathcal{L} : B^{\otimes \mathbb{N}} \to B^{\otimes \mathbb{N}}$ defined by

$$\blacktriangleright \ \alpha(f) = 1_B \otimes f$$

$$\blacktriangleright \mathcal{L}(f_1 \otimes f_2 \otimes ...) = \psi(f_1)f_2 \otimes f_3 \otimes ...$$

 $(B^{\otimes \mathbb{N}}, \alpha, \mathcal{L})$ is an Exel system.

Theorem (Brannan-Hamidi-I-Nelson-Wasilewski, 2022)

$$B^{\otimes \mathbb{N}} \rtimes_{\alpha, \mathcal{L}} \mathbb{N} \cong \mathcal{O}_{\dim B} \cong \mathcal{O}_E.$$

Exel systems for some quantum graphs

Example (Trivial quantum graph)

The quantum edge correspondence is $E \cong B$

 \rightsquigarrow quantum infinite path space should contain "infinite loops at each vertex."

Consider $E^{\otimes_B \mathbb{N}} = B \otimes_B B \otimes_B B \dots \cong B$ with $\alpha, \mathcal{L} : B^{\otimes \mathbb{N}} \to B^{\otimes \mathbb{N}}$ defined by

(B, id, id) is an Exel system.

Theorem (BHINW + Others)

$$B \rtimes_{id,id} \mathbb{N} \cong B \otimes C(\mathbb{T}) \cong \mathcal{O}_E.$$

Infinite path space for a quantum graph

In examples, the C^* -algebras on which we defined dynamics to build an Exel system came from $E^{\otimes_B \mathbb{N}}$. But in general, E is a C^* -correspondence over B, not a C^* -algebra, as is $E^{\otimes_B \mathbb{N}}$.

If $E^{\otimes_B \mathbb{N}}$ is not a C^* -algebra, we can't construct an Exel system in the same way as in the examples.

Thank you!